## On Pillai's problem with the Fibonacci and Pell sequences

Santos Hernández Hernández Florian Luca Luis Manuel Rivera

§1. Introduction. Let  $\mathbf{U} := (U_n)_{n \ge 0}$  and  $\mathbf{V} := (V_n)_{n \ge 0}$  be two linearly recurrent sequences of integers. Recently, the following variation of a problem of Pillai has been studied. Find all non-negative integer solutions  $(n, m, n_1, m_1)$  of the equation

$$U_n - V_m = U_{n_1} - V_{m_1}, \qquad (n,m) \neq (n_1, m_1).$$
 (1)

In particular, find also all integers c which can be written as the difference between an element of **U** and an element of **V** in at least two different ways. Pillai [12], studied this problem when **U** and **V** are the sequences of powers of a, and powers of b, respectively, where a, b are two given coprime integers different than  $0, \pm 1$ . It has been shown in [4] that, under some technical but natural conditions, equation (1) has only finitely many non-negative integer solutions and all of them are effectively computable. This version of Pillai's problem was initiated in [7] by Ddamulira, Luca and Rakotomalala who studied equation (1) when **U** and **V** are the sequences of Fibonacci numbers and powers of 2, respectively. Many other particular cases have been studied. See, for example [3], [6], [8]. We recall that the Fibonacci sequence  $(F_n)_{n\geq 0}$  is given by  $F_0 = 0$ ,  $F_1 = 1$  and the recurrence formula

$$F_{n+2} = F_{n+1} + F_n \qquad \text{for all} \qquad n \ge 0.$$

Let  $(P_n)_{n\geq 0}$  be the Pell sequence given by  $P_0 = 0$ ,  $P_1 = 1$ , and the recurrence formula

 $P_{n+2} = 2P_{n+1} + P_n \qquad \text{for all} \qquad n \ge 0.$ 

Their first terms are,

 $0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597, \ldots$ 

and

 $0, 1, 2, 5, 12, 29, 70, 169, 408, 985, 2378, 5741, 13860, 33461, \ldots,$ 

respectively. In this note, we study another particular case of this problem, namely equation (1) with Fibonacci and Pell numbers. More precisely, we look at the equation

$$F_n - P_m = F_{n_1} - P_{m_1} \tag{2}$$

in integer pairs  $(n,m) \neq (n_1,m_1)$ . Since  $F_1 = F_2 = 1$ , we assume that  $n \neq 1$ ,  $n_1 \neq 1$ . That is, whenever we think of 1 as a member of the Fibonacci sequence, we think of it as being  $F_2$ . Our result is then the following

**Theorem 1.** All solutions non-negative integer solutions  $(n, m, n_1, m_1)$  of (2) with  $n \neq 1$ ,  $n_1 \neq 1$  belong to the set

$$\left\{ \begin{array}{lll} (2,1,0,0), & (2,2,0,1), & (3,1,2,0), & (3,2,0,0), \\ (3,2,2,1), & (4,1,3,0), & (4,2,2,0), & (4,2,3,1), \\ (4,3,0,2), & (5,2,4,0), & (5,3,0,0), & (5,3,2,1), \\ (5,3,3,2), & (6,3,4,0), & (6,3,5,2), & (6,4,2,3), \\ (7,3,6,0), & (7,4,2,0), & (7,4,3,1), & (7,4,4,2), \\ (9,5,5,0), & (11,6,8,2), & (16,9,3,0), & (16,9,4,1) \end{array} \right\}.$$

The set of integers c admitting two representations as a difference between a Fibonacci and a Pell number in at least two different ways is

$$\{-4, -2, -1, 0, 1, 2, 3, 5, 8, 19\}$$

The representations of the above c are

$$\begin{array}{rcl} -4 &=& F_6 - P_4 = F_2 - P_3;\\ -2 &=& F_4 - P_3 = F_0 - P_2;\\ -1 &=& F_2 - P_2 = F_0 - P_1;\\ 0 &=& F_5 - P_3 = F_3 - P_2 = F_2 - P_1 = F_0 - P_0;\\ 1 &=& F_7 - P_4 = F_4 - P_2 = F_3 - P_1 = F_2 - P_0;\\ 2 &=& F_{16} - P_9 = F_4 - P_1 = F_3 - P_0;\\ 3 &=& F_6 - P_3 = F_5 - P_2 = F_4 - P_0;\\ 5 &=& F_9 - P_5 = F_5 - P_0;\\ 8 &=& F_7 - P_3 = F_6 - P_0;\\ 19 &=& F_{11} - P_6 = F_8 - P_2. \end{array}$$

**§2.** Tools. The first one is a lower bound for a linear forms in logarithms due to Matveev [11]. Let  $\alpha$  be an algebraic number of degree d. Let a be the leading coefficient of its minimal polynomial over  $\mathbb{Z}$  and let  $\alpha_1 = \alpha, \ldots, \alpha_d$  denote the conjugates of  $\alpha$ . The Weil height of  $\alpha$  is defined as

$$h(\alpha) = \frac{1}{d} \left( \log a + \sum_{i=1}^{d} \log \max\{|\alpha_i|, 1\} \right).$$

The height has the following basic properties. For  $\alpha, \beta$  algebraic numbers and  $m \in \mathbb{Z}$ , we have:

- $h(\alpha + \beta) \leq h(\alpha) + h(\beta) + \log 2$ .
- $h(\alpha\beta) \leq h(\alpha) + h(\beta)$ .
- $h(\alpha^m) = |m|h(\alpha)$ .

Now let  $\mathbb{L}$  be a real number field of degree  $d_{\mathbb{L}}, \alpha_1, \ldots, \alpha_\ell \in \mathbb{L}$  and  $b_1, \ldots, b_\ell \in \mathbb{Z} \setminus \{0\}$ . Let  $B \ge \max\{|b_1|, \ldots, |b_\ell|\}$  and

$$\Lambda = \alpha^{b_1} \cdots \alpha^{b_\ell} - 1.$$

Let  $A_1, \ldots, A_\ell$  be real numbers such that

$$A_i \ge \max\{d_{\mathbb{L}}h(\alpha_i), |\log \alpha_i|, 0.16\}$$
 for all  $i = 1, \dots, \ell$ .

The following result is due to Matveev in [11] (see also Theorem 9.4 in [2]).

**Theorem 2.** Assume that  $\Lambda \neq 0$ . Then

$$\log |\Lambda| > -1.4 \times 30^{\ell+3} \times \ell^{4.5} \times d_{\mathbb{L}}^2 (1 + \log d_{\mathbb{L}}) (1 + \log B) A_1 \cdots A_\ell.$$

In this paper, we always use  $\ell = 3$ . Further,  $\mathbb{L} = \mathbb{Q}[\sqrt{2}, \sqrt{5}]$  has degree  $d_{\mathbb{L}} = 4$ . Thus, once for all we fix the constant

$$C := 5.46696 \times 10^{12} > 1.4 \times 30^{3+3} \times 3^{4.5} \times 4^2(1 + \log 4).$$

Matveev's bound gives us some large bounds on our parameters. In order to lower such bounds, we use a version of a reduction method of Baker-Davenport based on Lemma in [1]. We shall use the one given by Bravo, Gomez and Luca in [5]. For a real number x, we write

$$||x|| = \min\{|x - n| : n \in \mathbb{Z}\}.$$

**Lemma 3.** Let M be a positive integer. Let  $\tau$ ,  $\mu$ , A > 0, B > 1 be given real numbers. Assume that p/q is a convergent of  $\tau$  such that q > 6M and  $\varepsilon := ||q\mu|| - M||q\tau|| > 0$ . Then the inequality

$$0 < |n\tau - m + \mu| < \frac{A}{B^w}$$

does not have a solution in positive integers n, m and w in the ranges

$$n \leqslant M$$
 and  $w \geqslant \frac{\log (Aq/\varepsilon)}{\log B}$ .

This lemma is a slightly variation of the one given by Dujella and Petho in [9]. The following lemma is also useful. It is Lemma 7 in [10].

**Lemma 4.** If  $m \ge 1$ ,  $T > (4m^2)^m$  and  $T > x/(\log x)^m$ , then

 $x < 2^m T (\log T)^m.$ 

**§3. Proof of Theorem 1.** We start with some basic properties of our sequences. Put

$$\alpha := \frac{1+\sqrt{5}}{2}, \quad \beta := \frac{1-\sqrt{5}}{2}; \quad \text{and} \quad \gamma := 1+\sqrt{2}, \quad \delta := 1-\sqrt{2}.$$

We have the well-known Binet's formulas

$$F_n = \frac{\alpha^n - \beta^n}{\sqrt{5}}$$
 and  $P_n = \frac{\gamma^n - \delta^n}{2\sqrt{2}}$  (3)

which hold for all  $n \ge 0$ . Further, the inequalities

$$\alpha^{n-2} \leqslant F_n \leqslant \alpha^{n-1} \quad \text{and} \quad \gamma^{n-2} \leqslant P_n \leqslant \gamma^{n-1}$$
(4)

also hold for all  $n \ge 1$ .

Now, we study our equation (2) in non-negative integers  $(n, m, n_1, m_1)$  with  $(n, m) \neq (n_1, m_1)$ . As we said, we assume  $n \neq 1$ ,  $n_1 \neq 1$ . It could happen that  $\min\{n, n_1\} = 0$ . At any rate,  $\max\{n, n_1\} \ge 2$ . If in (2) we have  $m = m_1$ , then  $F_n = F_{n_1}$ , implies that  $n = n_1$ , a contradiction. Thus, from now on we assume  $m > m_1$ . Rewriting (2) as

$$F_n - F_{n_1} = P_m - P_{m_1}, (5)$$

we observe the right-hand side is positive. Hence, so is the left-hand side, therefore  $n > n_1$ . We now compare the two sides of (5) using (4). We have

$$\alpha^{n-4} \leqslant F_n - F_{n_1} = P_m - P_{m_1} \leqslant P_m \leqslant \gamma^{m-1}.$$

The left-hand side inequality is clear if  $n_1 = 0$ . It is also clear if  $n_1 \neq 0$ , since in that case  $n_1 \geq 2$ , so  $n \geq 3$ , so  $F_n - F_{n_1} \geq F_n - F_{n-1} = F_{n-2} \geq \alpha^{n-4}$ . Thus,  $\alpha^{n-4} \leq \gamma^{m-1}$ . In a similar way,

$$\alpha^{n-1} \geqslant F_n \geqslant F_n - F_{n_1} = P_m - P_{m_1} \geqslant P_{m-1} \geqslant \gamma^{m-3},$$

where the right-most inequality is clear (both for  $m_1 = 0$  and for  $m_1 > 0$ ). We thus have

$$n-4 \leq (m-1)\frac{\log \gamma}{\log \alpha}$$
 and  $n-1 \geq \frac{\log \gamma}{\log \alpha}(m-3).$  (6)

Since  $\log \gamma / \log \alpha = 1.8315709239...$  it follows that if  $n \leq 300$ , then  $m \leq 167$ . Running a Mathematica program in the range  $0 \leq n_1 < n \leq 300$  and  $0 \leq n_1 < n \leq 300$  $m_1 < m \leq 167$ , with our convention, we obtain all the possibilities listed in Theorem 1.

From now on, n > 300. Further, by (6) we get m > 163 and also n > m. From Binet's formulas (3), we obtain

$$\left| \frac{\alpha^{n}}{\sqrt{5}} - \frac{\gamma^{m}}{2\sqrt{2}} \right| = \left| \frac{\alpha^{n_{1}} + \beta^{n} - \beta^{n_{1}}}{\sqrt{5}} - \frac{\gamma^{m_{1}} - \delta^{m_{1}} + \delta^{m}}{2\sqrt{2}} \right| \leqslant \frac{\alpha^{n_{1}} + 2}{\sqrt{5}} + \frac{\gamma^{m_{1}} + 2}{2\sqrt{2}} \\ \leqslant 2 \max\{\alpha^{n_{1}+2}, \gamma^{m_{1}+1}\}.$$

$$(7)$$

Dividing through by  $\gamma^m/2\sqrt{2}$  we get

$$\left|\frac{4}{\sqrt{10}}\gamma^{-m}\alpha^n - 1\right| \leqslant \max\{\alpha^{n_1-n+9}, \gamma^{m_1-m+1}\},\tag{8}$$

where we have used that  $\alpha^{n-4} \leq \gamma^{m-1}$  as well as the fact that  $4\sqrt{2} < \lambda^2 < \alpha^4$ . Let  $\Lambda$  be the expression inside the absolute value in the left-hand side above. Observe that  $\Lambda$  is not zero. Indeed, otherwise  $8/5 = \gamma^{2m}/\alpha^{2n}$  is both a unit (an algebraic integer whose reciprocal is also an algebraic integer) and a rational number, which is false since the only rational units are  $\pm 1$ .

Now we apply Matveev's inequality with

$$\alpha_1 = \frac{4}{\sqrt{10}}, \quad \alpha_2 = \gamma, \quad \alpha_3 = \alpha, \quad b_1 = 1, \quad b_2 = -m, \quad b_3 = n.$$

We have B = n. Further, we have  $h(\alpha_1) = (\log 8)/2$ ,  $h(\alpha_2) = (\log \gamma)/2$  and  $h(\alpha_3) = \log \alpha/2$ . Thus, we may take  $A_1 := 4.2, A_2 := 1.8$  and  $A_3 := 1$  we obtain that

 $\log |\Lambda| > -C(1 + \log n) \times 4.2 \times 1.8.$ 

Comparing with (8) we obtain

$$\min\{(n - n_1 - 9)\log\alpha, (m - m_1 - 1)\log\gamma\} \leq 4.13302 \times 10^{13}(1 + \log n).$$
(9)

We next study each of these two possibilities.

**Case 1.**  $\min\{(n - n_1) \log \alpha, (m - m_1) \log \gamma\} = (n - n_1) \log \alpha.$ 

To this case, we rewrite our equation as follows:

$$\left| \left( \frac{\alpha^{n-n_1} - 1}{\sqrt{5}} \right) \alpha^{n_1} - \frac{\gamma^m}{2\sqrt{2}} \right| = \left| \frac{\beta^n - \beta^{n_1}}{\sqrt{5}} - \frac{\gamma^{m_1} - \delta^{m_1} + \delta^m}{2\sqrt{2}} \right|$$
  
 
$$\leqslant \frac{2}{\sqrt{5}} + \frac{\gamma^{m_1} + 2}{2\sqrt{2}} < \gamma^{m_1+2}.$$

Thus,

$$\left(\frac{4(\alpha^{n-n_1}-1)}{\sqrt{10}}\right)\alpha^{n_1}\gamma^{-m}-1\Big|<\gamma^{m_1-m+4}.$$
(10)

Let  $\Lambda_1$  be the expression inside the absolute value which is in the left-hand side. We note that  $\Lambda_1 \neq 0$ , for if this is not so then we would get

$$\frac{\alpha^n - \alpha^{n_1}}{\gamma^m} = \frac{\sqrt{10}}{4},$$

which implies that the right-hand side is an algebraic integer, which it isn't (it's square is 5/8). We apply again Matveev's inequality by taking

$$\alpha_1 = \frac{4(\alpha^{n-n_1} - 1)}{\sqrt{10}}, \quad \alpha_2 = \gamma, \quad \alpha_3 = \alpha, \quad b_1 = 1, \quad b_2 = -m, \quad b_3 = n_1.$$

Thus, B = n. The heights of  $\alpha_2$  and  $\alpha_3$  have already been calculated. As for  $h(\alpha_1)$ , we have

$$h\left(\frac{4(\alpha^{n-n_1}-1)}{\sqrt{10}}\right) \leqslant h\left(\frac{4}{\sqrt{10}}\right) + h\left(\alpha^{n-n_1}-1\right) \leqslant \frac{\log 8}{2} + h(\alpha^{n-n_1}) + \log 2$$
$$= \frac{\log 32}{2} + (n-n_1)\frac{\log \alpha}{2} \leqslant \frac{4.13304 \times 10^{13}(1+\log n)}{2},$$

where we have used (9). Thus, we can take  $A_1 := 8.26608 \times 10^{13} (1 + \log n)$ ,  $A_2$  and  $A_3$  as in the analysis of  $\Lambda$ , and get

$$\log |\Lambda_1| > -C \times (8.26608 \times 10^{13} (1 + \log n)^2) \times 1.8.$$

Combining this with (10), we get

$$(m - m_1)\log\gamma < 8.13424 \times 10^{26} (1 + \log n)^2.$$

Case 2.  $\min\{(n - n_1)\log \alpha, (m - m_1)\log \gamma\} = (m - m_1)\log \gamma.$ Here, we rewrite our equation as

Here, we rewrite our equation as

$$\left| \frac{\alpha^{n}}{\sqrt{5}} - \left( \frac{\gamma^{m-m_{1}} - 1}{2\sqrt{2}} \right) \gamma^{m_{1}} \right| = \left| \frac{\beta^{n} + \alpha^{n_{1}} - \beta^{n_{1}}}{\sqrt{5}} - \frac{\delta^{m} - \delta^{m_{1}}}{2\sqrt{2}} \right|$$
$$\leqslant \frac{\alpha^{n_{1}} + 2}{\sqrt{5}} + \frac{1}{\sqrt{2}} < \alpha^{n_{1}+5}.$$

Thus,

$$\left| 1 - \left( \frac{\sqrt{10}(\gamma^{m-m_1} - 1)}{4} \right) \gamma^{m_1} \alpha^{-n} \right| < \alpha^{n_1 - n + 7}.$$
 (11)

We let  $\Lambda_2$  be the expression inside the absolute value in the left-hand side. As before,  $\Lambda_2 \neq 0$ , for otherwise we get that 8/5 is an algebraic integer, which is false. We apply again Matveev's inequality by taking

$$\alpha_1 = \frac{\sqrt{10}(\gamma^{m-m_1} - 1)}{4}, \quad \alpha_2 = \gamma, \quad \alpha_3 = \alpha, \quad b_1 = 1, \quad b_2 = m_1, \quad b_3 = -n.$$

Thus, B = n. The heights of  $\alpha_2$  and  $\alpha_3$  have already been calculated. As for  $h(\alpha_1)$ , we have

$$h\left(\frac{\sqrt{10}(\gamma^{m-m_1}-1)}{4}\right) \leqslant h\left(\frac{\sqrt{10}}{4}\right) + h\left(\gamma^{m-m_1}-1\right)$$
$$\leqslant \frac{4.13304 \times 10^{13}(1+\log n)}{2},$$

Thus, we can take the same  $A_1$  as in Case 1, and so we get the same lower bound for  $\log |\Lambda_2|$ . Therefore,

$$(n - n_1)\log\gamma < 8.13424 \times 10^{26} (1 + \log n)^2.$$

So, we have proved that

$$\max\{(n-n_1)\log\alpha, (m-m_1)\log\gamma\} \leqslant 8.13424 \times 10^{26}(1+\log n)^2.$$
 (12)

We now get a bound on n. Using Binet's formulas (3), we write our equation as follows:

$$\left|\frac{\alpha^{n-n_1}-1}{\sqrt{5}}\alpha^{n_1}-\frac{\gamma^{m-m_1}-1}{2\sqrt{2}}\gamma^{m_1}\right| = \left|\frac{\beta^n-\beta^{n_1}}{\sqrt{5}}-\frac{\delta^m-\delta^{m_1}}{2\sqrt{2}}\right| < \frac{2}{\sqrt{5}} + \frac{1}{\sqrt{2}} < 2.$$

Dividing across by  $(\gamma^m - \gamma^{m-1})/2\sqrt{2}$ , we obtain

$$\left| \left( \frac{4}{\sqrt{10}} \left( \frac{\alpha^{n-n_1} - 1}{\gamma^{m-m_1} - 1} \right) \right) \gamma^{-m_1} \alpha^{n_1} - 1 \right| < \frac{4\sqrt{2}}{\gamma^m - \gamma^{m_1}} < \frac{8\sqrt{2}}{\gamma^m} < \frac{1}{\alpha^{n-8}}, \quad (13)$$

where we used  $\alpha^{n-4} < \gamma^{m-1}$ , as well as the fact that  $8\sqrt{2} < \alpha^4 \gamma$ . We let  $\Lambda_3$  be the expression inside the absolute value in (13). We apply Matveev's inequality with

$$\alpha_1 = \frac{4}{\sqrt{10}} \left( \frac{\alpha^{n-n_1} - 1}{\gamma^{m-m_1} - 1} \right), \ \alpha_2 = \gamma, \ \alpha_3 = \alpha, \ b_1 = 1, \ b_2 = -m_1, \ b_3 = -n_1.$$

Thus, we take B = n. We need to show that  $\Lambda_3 \neq 0$ . To do this we take the  $\mathbb{Q}$ -automorphism  $\sigma$  of  $\mathbb{L}$  given by  $\sigma(\sqrt{5}) = -\sqrt{5}$  and  $\sigma(\sqrt{2}) = \sqrt{2}$ . Under this

automorphism, we have  $\sigma(\alpha) = \beta$ ,  $\sigma(\gamma) = \gamma$  and  $\sigma(\sqrt{10}) = -\sqrt{10}$ . Thus, if  $\Lambda_3 = 0$ , then  $\sigma(L_3) = 0$ , which implies, in particular, that

$$\frac{\sqrt{10}}{4} = \left| \frac{\beta^n - \beta^{n_1}}{\gamma^m - \gamma^{m_1}} \right| < \frac{2}{\gamma^m (\gamma - 1)} < \frac{1}{2},$$

since m > 163, which is a contradiction. As before, the heights of  $\alpha_2$  and  $\alpha_3$  have already been calculated. For  $h(\alpha_1)$ , we have

$$h\left(\frac{4}{\sqrt{10}}\left(\frac{\alpha^{n-n_1}-1}{\gamma^{m-m_1}-1}\right)\right) \leqslant h\left(\frac{4}{\sqrt{10}}\right) + h\left(\alpha^{n-n_1}+1\right) + h\left(\gamma^{m-m_1}+1\right)$$
$$\leqslant \frac{\log 128}{2} + (n-n_1)\frac{\log \alpha}{2} + (m-m_1)\frac{\log \gamma}{2}$$
$$\leqslant 8.13425 \times 10^{26}(1+\log n)^2.$$

Thus, we can take  $A_1 := 3.25368 \times 10^{27} (1 + \log n)^2$ , and  $A_2, A_3$  as before. Therefore, we get

$$\log |\Lambda_3| > -C(1 + \log n) \times (3.25368 \times 10^{27} (1 + \log n)^2) \times 1.8$$
  
> -3.20181 \times 10^{40} (1 + \log n)^3,

which, upon comparing it to (13) and applying Lemma 4, we obtain

$$n < 3.77669 \times 10^{48}. \tag{14}$$

Now, we will reduce the upper bound of n. To do this, let  $\Gamma$  be defined as

$$\Gamma = n \log \alpha - m \log \gamma + \log \left(\frac{4}{\sqrt{10}}\right).$$

Assume first that  $\min\{n - n_1, m - m_1\} \ge 20$ . We note that  $\Lambda = e^{\Gamma} - 1 \ne 0$ , so  $\Gamma \ne 0$ . If  $\Gamma > 0$  then

$$0 < \Gamma < e^{\Gamma} - 1 = \Lambda = |\Lambda| < \max\{\alpha^{n_1 - n + 9}, \gamma^{m_1 - m + 1}\}.$$

On the other hand, if  $\Gamma < 0$ , we then have  $1 - e^{\Gamma} = |e^{\Gamma} - 1| < 1/2$  which implies  $e^{|\Gamma|} < 2$ . Thus,

$$0 < |\Gamma| < e^{|\Gamma|} - 1 = e^{|\Gamma|} |\Lambda| < 2 \max\{\alpha^{n_1 - n + 9}, \gamma^{m_1 - m + 1}\}.$$

So, in both cases we have

$$0 < |\Gamma| < 2\max\{\alpha^{n_1 - n + 9}, \gamma^{m_1 - m + 1}\}.$$
(15)

Dividing through by  $\log \gamma$  in the above inequality, we get

$$0 < |n\tau - m + \mu| < \max\left\{\frac{175}{\alpha^{n-n_1}}, \frac{6}{\gamma^{m-m_1}}\right\},$$

where

$$\tau := \frac{\log \alpha}{\log \gamma}, \qquad \mu := \frac{\log \left(4/\sqrt{10}\right)}{\log \gamma}.$$

Now we apply Lemma 3. To do this, we take  $M := 3.77669 \times 10^{48}$  (a bound on *m* and *n* by (14)) our  $\tau$  and, with a *Mathematica* program, we find that the denominator of the convergent

$$\frac{p_{112}}{q_{112}} = \frac{111842821415068814601069451383096958405345992106163812}{204848059751598401563305907296432335323118859258712413}$$

of  $\tau$  satisfies  $q_{112} > 6M$  and that  $\varepsilon = ||q\mu|| - M||q\tau|| = 0.105822 > 0$ . This implies, with  $(A, B) = (175, \alpha)$  or  $(6, \gamma)$ , that either

$$n - n_1 \leq 271$$
, or  $m - m_1 \leq 144$ .

We now look at each one of these two cases. First, we assume that  $n - n_1 \leq 271$ and  $m - m_1 \geq 20$ . In this case, we consider

$$\Gamma_1 = n_1 \log \alpha - m \log \gamma + \log \left(\frac{4(\alpha^{n-n_1} - 1)}{\sqrt{10}}\right).$$

As before,  $e^{\Gamma_1} - 1 = \Lambda_1 \neq 0$ , so  $\Gamma_1 \neq 0$ . We go to (10). With an argument similar to a previous one, we have that

$$0 < |\Gamma_1| < \frac{2\gamma^4}{\gamma^{m-m_1}}.$$

Dividing through by  $\log \gamma$  we obtain

$$0 < |n_1 \tau - m + \mu| < \frac{78}{\gamma^{m - m_1}},$$

where  $\tau$  is the same one as above and

$$\mu := \frac{\log\left(4(\alpha^{n-n_1} - 1)/\sqrt{10}\right)}{\log\gamma}.$$

We apply again Lemma 3 noting that  $n_1 > 0$ , for otherwise we would have that  $n \leq 271$  which contradicts our hypothesis that n > 300. Consider

$$\mu_k := \frac{\log(4(\alpha^k - 1)/\sqrt{10})}{\log \gamma}, \quad \text{for} \quad k = 1, \dots, 271.$$

We ran a *Mathematica* program and found that the same convergent  $p_{112}/q_{112}$  satisfies  $q_{112} > 6M$ . Further,  $\varepsilon_k \ge 0.00119532$  for all  $1 \le k \le 271$ . For each of

the values of  $\varepsilon_k$  and with  $(A, B) = (78, \gamma)$ , we calculate  $\log (78q_{112}/\varepsilon_k) / \log \gamma$ and found that each of them is at most 152. Thus,  $m - m_1 \leq 152$ .

Now let us look at the other case. Assume that  $m - m_1 \leq 144$  and  $n - n_1 \geq 20$ . We consider

$$\Gamma_2 = n \log \alpha - m_1 \log \gamma + \log \left(\frac{4}{\sqrt{10}(\gamma^{m-m_1} - 1)}\right).$$

We note that  $1 - e^{-\Gamma_2} = \Lambda_2 \neq 0$ , so  $\Gamma_2 \neq 0$ . We go to (11). With an argument similar to one above, we obtain

$$0 < |\Gamma_2| < \frac{2\alpha^7}{\alpha^{n-n_1}}.$$

Dividing through by  $\log \lambda$ , we get

$$0 < |n\tau - m_1 + \mu| < \frac{66}{\alpha^{n-n_1}},$$

where  $\tau$  is the same one as above and

$$\mu := \frac{\log\left(4/(\sqrt{10}(\gamma^{m-m_1}-1))\right)}{\log\gamma}$$

Now we use again Lemma 3 noting that  $m_1 > 0$ , which is the case, since otherwise we have  $m \leq 144$ , which contradicts our hypothesis m > 163. As above, by considering now

$$\mu_{\ell} := \frac{\log\left(4/(\sqrt{10}(\gamma^{\ell} - 1))\right)}{\log \gamma}, \quad \text{for all} \quad \ell = 1, \dots, 144$$

and running a *Mathematica* program, we find that  $q_{112} > 6M$ , and that for this convergent  $\varepsilon_{\ell} \ge 0.0000620747$  for all  $1 \le \ell \le 144$ . For each of these  $\varepsilon_{\ell}$ and with  $(A, B) := (66, \alpha)$ , we calculated  $\log (66q_{112}/\varepsilon_{\ell}) / \log \alpha$  and found that all these numbers are at most 156. Thus  $n - n_1 \le 156$ .

So, we got that either  $n - n_1 \leq 271$  or  $m - m_1 \leq 144$ . Assuming the first one we deduced  $m - m_1 \leq 152$ , and assuming the second one, we deduced  $n - n_1 \leq 156$ . Altogether, we have  $n - n_1 \leq 271$ ,  $m - m_1 \leq 152$ . So, it remains to study this case. We consider

$$\Gamma_3 = n_1 \log \alpha - m_1 \log \gamma + \log \left(\frac{4}{\sqrt{10}} \left(\frac{\alpha^{n-n_1} - 1}{\gamma^{m-m_1} - 1}\right)\right).$$

We note that  $e^{\Gamma_3} - 1 = \Lambda_3$ . Again, since  $\Lambda_3 \neq 0$ , we have that  $\Gamma_3 \neq 0$ . Since n > 300, we get

$$0 < |\Gamma_3| < \frac{2\alpha^8}{\alpha^n}.$$

Dividing through by  $\log \gamma$ , we get

$$0 < |n_1 \tau - m_1 + \mu| < \frac{107}{\alpha^n},$$

where  $\tau$  is as above and

$$\mu := \frac{\log \left( 4(\alpha^{n-n_1} - 1) / \sqrt{10}(\gamma^{m-m_1} - 1) \right)}{\log \gamma}.$$

We apply for the last time Lemma (3). As above, we have that  $n_1, m_1 > 0$ . Thus, we consider

$$\mu_{k,\ell} := \frac{\log\left(4(\alpha^k - 1)/\sqrt{10}(\gamma^\ell - 1)\right)}{\log \gamma}, \quad k = 1, \dots, 271, \quad \ell = 1, \dots, 152.$$

Running a Mathematica program, we find again that the same convergent works namely  $q_{112} > 6M$  and  $\varepsilon_{k,\ell} \ge 0.0000307768$  for all  $1 \le k \le 271$  and  $1 \le \ell \le 152$ . For each of these values  $\varepsilon_{k,\ell}$ , with  $(A, B) := (107, \alpha)$ , we calculated log  $(107q_{112}/\varepsilon_{k,\ell}) / \log \alpha$  and found that the maximum value of them is  $\le 157$ . Thus,  $n \le 157$ , which contradicts our assumption on n. This finishes the proof of our theorem.

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Santos HERNÁNDEZ HERNÁNDEZ Luis Manuel RIVERA Unidad Académica de Matemáticas Universidad Autónoma de Zacatecas Calzada Solidaridad esquina Camino a la Bufa S/N C.P. 98000 Zacatecas, Zac. MEXICO *E-mail:* shh@matematicas.reduaz.mx, luismanuel.rivera@gmail.com

Florian LUCA School of Mathematics University of the Witwatersrand Private Bag X3, WITS 2050 Johannesburg, SOUTH AFRICA; Max Planck Institute for Mathematics, Vivatsgasse 7, 53111 Bonn, GERMANY; Department of Mathematics, Faculty of Sciences, University of Ostrava, 30 Dubna 22, 701 03 Ostrava 1, CZECH REPUBLIC *E-mail:* Florian.Luca@wits.ac.za